

Automata

course based on Jurafsky and Martin [2009, Chap.2]



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Presentation Plan

- 1 **Automata**
 - Definition
 - Exercise 1 on Formalism
 - Exercise 2 on Regex

- 2 **Determinism and Automata**
 - Definition
 - Example

- 3 **Summary**

Table of Contents

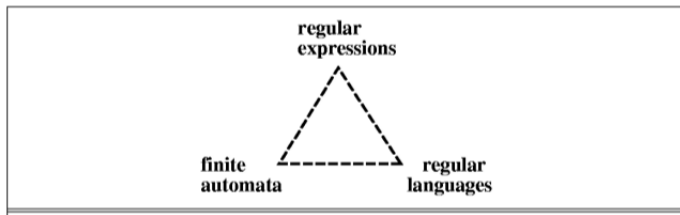
- 1 **Automata**
 - Definition
 - Exercise 1 on Formalism
 - Exercise 2 on Regex
- 2 **Determinism and Automata**
 - Definition
 - Example
- 3 **Summary**

Table of Contents

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 - Definition
 - Exercise 1 on Formalism
 - Exercise 2 on Regex
- 2 **Determinism and Automata**
 - Definition
 - Example
- 3 **Summary**

Why Automata ?

Every regular expression are based on a formalism called automata. Without knowing it, you apply it all the time. It is a basic formalism on which you will build your knowledge to understand other formalisms.



Definition

Finite-state automaton (FSA, plural : automata), or simply a state machine, is a mathematical model of computation used to design computer programs and regular expressions. It is conceived as an abstract machine that can be in one of a finite number of **states**.

The machine is in only one state at a time ; the state it is in at any given time is called the **current state**. It can change from one state to another when initiated by a triggering event or condition ; this is called a **transition**.

A particular FSM is defined by a list of its states, and the triggering condition for each transition. **Deterministic finite automaton** (DFA) and **nondeterministic finite automaton** (NFA) are subcategories of Finite-state automaton

Definition

An automaton can be described by its formal definition. Here is the lexicon of FSA formalism :

- Q : a finite set of N states q_0, q_1, \dots, q_N
- Σ : a finite input alphabet of symbols
- q_0 : the start state
- F : the set of final states, $F \subseteq Q$
- $\delta(q, i)$: the transition function or transition matrix between states. Given a state $q \in Q$ and an input symbol $i \in \Sigma$, $\delta(q, i)$ returns a new state $q' \in Q$. δ is thus a relation from $Q \times \Sigma$ to Q ;

Definition

An automaton can as well get a visual representation

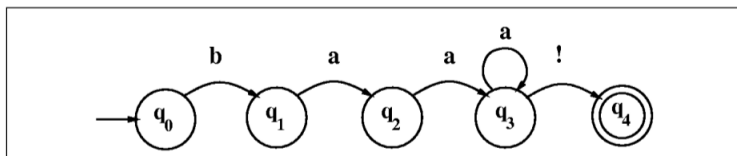


Figure 2.10 A finite-state automaton for talking sheep.

For the sheeptalk automaton in Figure 2.10, $Q = \{q_0, q_1, q_2, q_3, q_4\}$,
 $\Sigma = \{a, b, !\}$, $F = \{q_4\}$

Table of Contents

- 1 Automata**
 - Definition
 - Exercise 1 on Formalism
 - Exercise 2 on Regex
- 2 Determinism and Automata**
 - Definition
 - Example
- 3 Summary**

Exercise 1 on Formalism

The next slide describes an Automaton

According to the formalism, color/circle what represents :

- Q in blue
 - Σ in green
 - q_0 and F (if not missing) in dotted blue
 - $\delta(q,i)$ in red
- Q : a finite set of N states q_0, q_1, \dots, q_N
 - Σ : a finite input alphabet of symbols
 - q_0 : the start state
 - F : the set of final states, $F \subseteq Q$
 - $\delta(q,i)$: the transition function or transition matrix between states. Given a state $q \in Q$ and an input symbol $i \in \Sigma$, $\delta(q,i)$ returns a new state $q' \in Q$. δ is thus a relation from $Q \times \Sigma$ to Q ;

Exercise

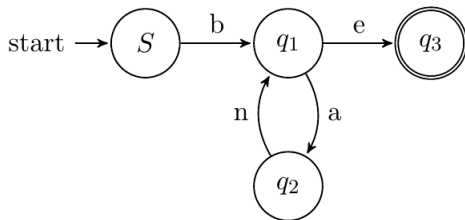


Figure 1: Finite State Automaton, accepting the pattern $b(an)^+e$

Table of Contents

- 1 Automata**
 - Definition
 - Exercise 1 on Formalism
 - Exercise 2 on Regex

- 2 Determinism and Automata**
 - Definition
 - Example

- 3 Summary**

Exercise 2 on Regex

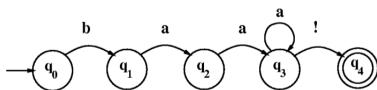


Figure 2.10 A finite-state automaton for talking sheep.

Instructions 1

ba

baaa !

baaa

baa !

b

baaaaaaaaaaaaaa !

- Which of those examples should be accepted by this automaton ?

Exercise 2

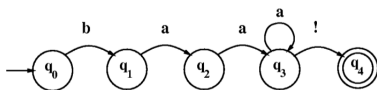


Figure 2.10 A finite-state automaton for talking sheep.

Instructions

ba

baaa !

baaa

baa !

b

baaaaaaaaaaaaaa !

- Open a terminal window and find the grep regular expression that correspond to the image up. Apply it to the file, do you get the right output ?

Exercise 2 on Regex

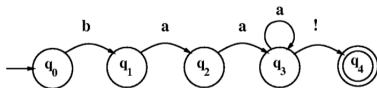


Figure 2.10 A finite-state automaton for talking sheep.

Conclusion on the exercises

Congratulations you went all the way from formalism to implementation !

Table of Contents

- 1 **Automata**
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 - Exercise 1 on Formalism
 - Exercise 2 on Regex

- 2 **Determinism and Automata**
 - Definition
 - Example

- 3 **Summary**

Table of Contents

- 1 **Automata**
 - Definition
 - Exercise 1 on Formalism
 - Exercise 2 on Regex
- 2 **Determinism and Automata**
 - Definition
 - Example
- 3 **Summary**

Definition

In theory of computation a **deterministic finite automaton (DFA)**—also known as deterministic finite acceptor (DFA) is a finite state machine that accepts/rejects finite strings of symbols and which transitions is determined only by the current state and the symbol it is looking at.

If in an automaton a state can lead to two ways with the same symbol this automaton is called non-deterministic finite automaton.

Table of Contents

- 1 **Automata**
 - Definition
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 - Exercise 2 on Regex
- 2 **Determinism and Automata**
 - Definition
 - Example
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Example Deterministic VS Non-deterministic

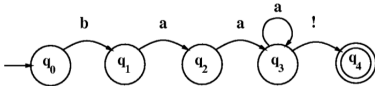


Figure 2.10 A finite-state automaton for talking sheep.

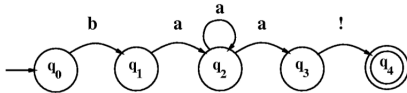


Figure 2.18 A non-deterministic finite-state automaton for talking sheep (NFA #1). Compare with the deterministic automaton in Figure 2.10.

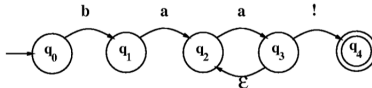


Figure 2.19 Another NFA for the sheep language (NFA #2). It differs from NFA #1 in Figure 2.18 in having an ϵ -transition.

The non-determinism need to be handled by algorithms like Back-up, look-ahead and parallelism in order avoid errors.

Table of Contents

- 1 **Automata**
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Summary

- We saw a formalism called **automaton**
- Automaton is the basis of several other formalism and of Regex
- Automata can be **deterministic** or **non-deterministic**
- They are a way to represent the step of a reasoning in the most computational friendly way.
- To be academic, you must always check that all the element of the formalism is present in the automaton representation you create (Start, Final states, symboles etc.)

Daniel Jurafsky and James H Martin. *Speech and Language Processing : An Introduction to Natural Language Processing, Computational Linguistics, and Speech Recognition*, volume 163 of *Prentice Hall Series in Artificial Intelligence*. Prentice Hall, 2009.